

METIS: A Non-Clairvoyant, Workflow-Aware OS Scheduler for Serverless Applications

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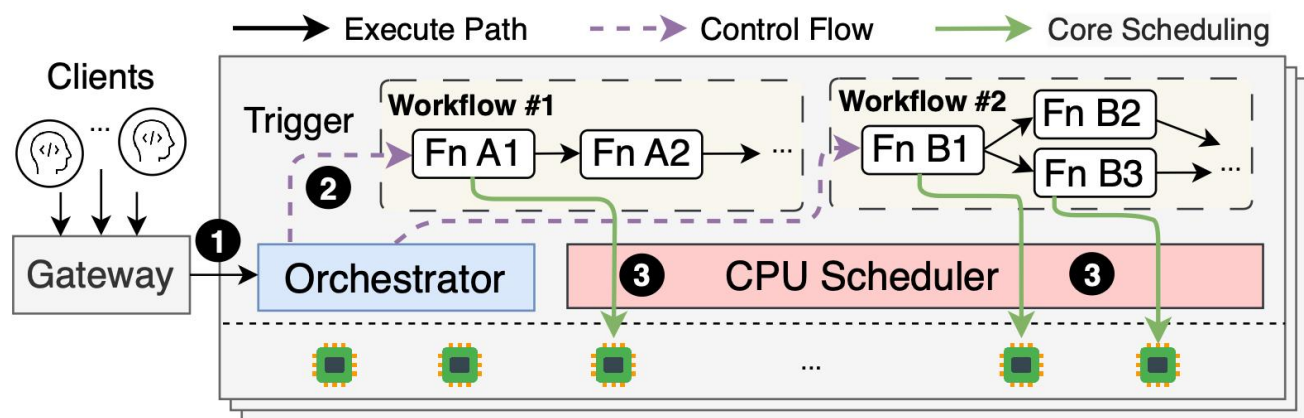
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Serverless Workflow Execution Model

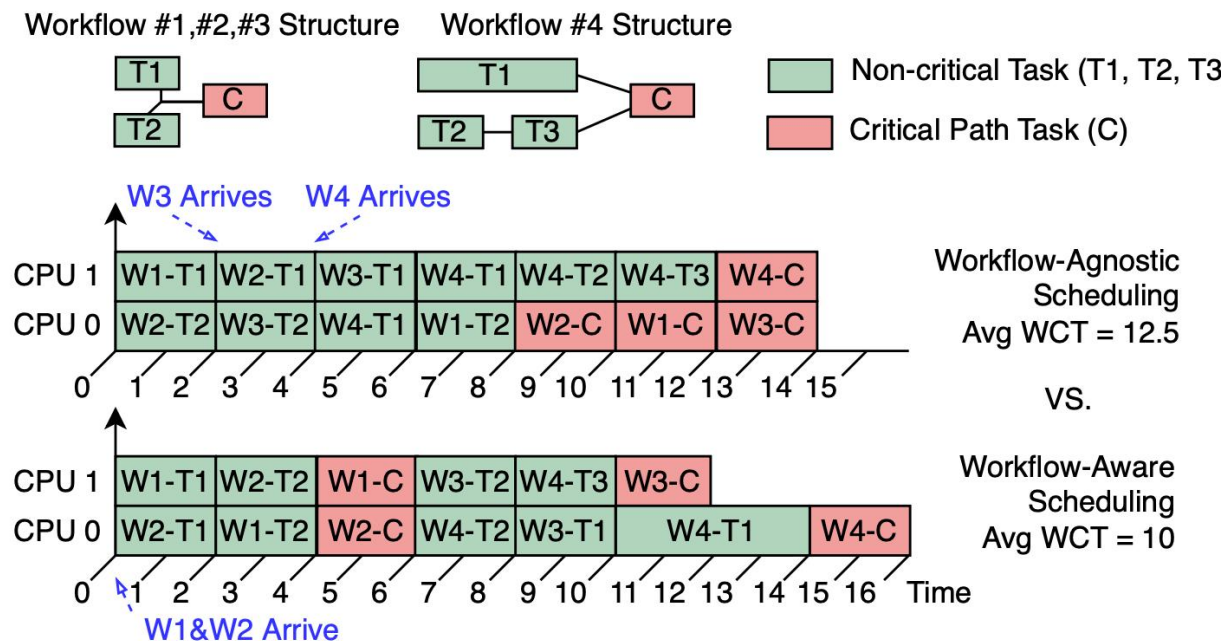


• Host OS Scheduler

- Each function runs as an independent task (container/microVM) enqueued in the OS runqueue.
- Linux CFS (typical) allocates CPU **without workflow** context.

- **Gateway**: Handles auth, validation, routing; forwards triggers to the workflow orchestrator.
- **Orchestrator**: Manages a DAG: nodes = functions, edges = data/control dependencies.
 - Dispatches a function only when dependencies are satisfied; exploits parallel branches.

Workflow-OS Scheduling Mismatch (Case Study)



Comparison of workflow-agnostic (up) and workflow-aware (down) scheduling for 4 concurrent workflows on a dual-core system.

- Workflow-agnostic scheduling:
 - Schedules functions independently, without workflow context
 - Intra-workflow blocking:** Critical-path tasks can be delayed by non-critical tasks **within the same workflow**
 - Inter-workflow blocking:** Critical tasks may also be blocked by functions from **other workflows** competing for CPU
 - Increased workflow completion time (WCT)
- Workflow-aware scheduling:
 - Prioritizes critical-path tasks within workflows
 - Eliminates unnecessary blocking & reduces overall workflow latency

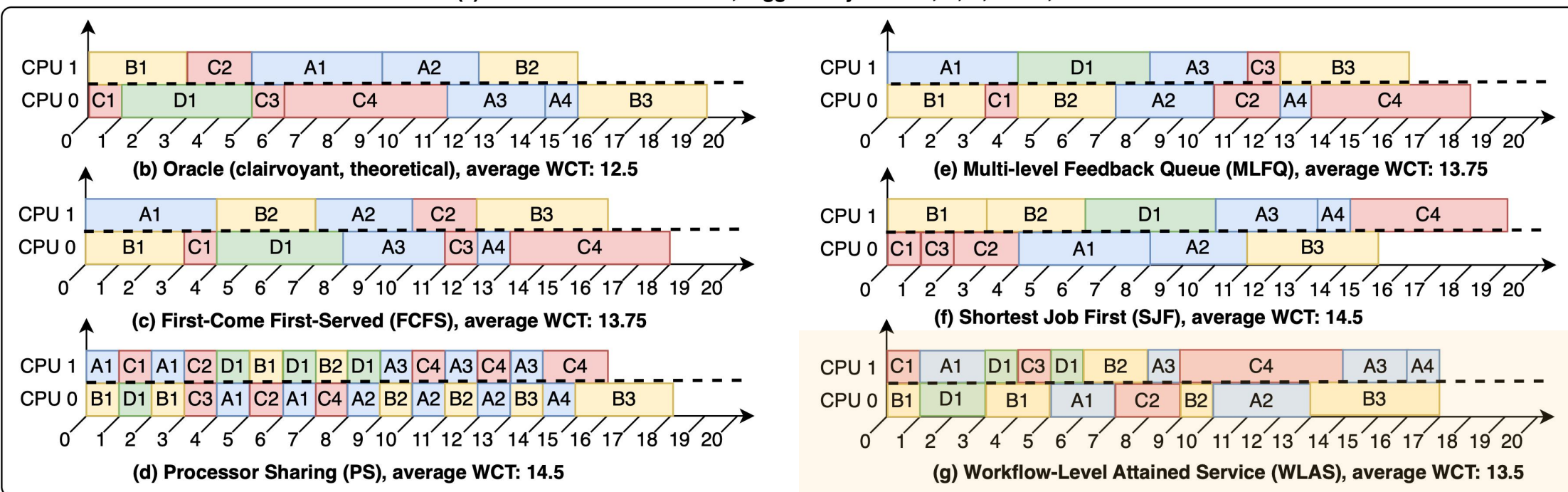
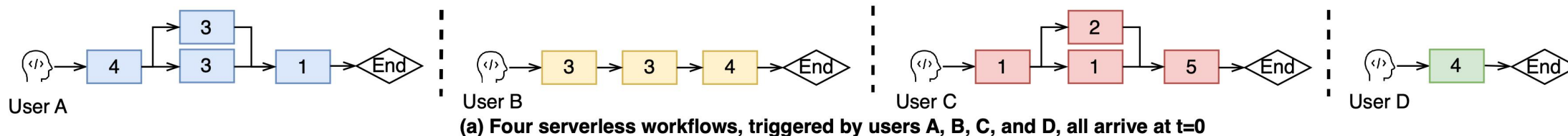
Related Work

- Existing research focuses on:
 - Cold start latency reduction [SONIC(ATC'21), AlloyStack(EuroSys'25)...]
 - Elastic resource scaling [Jiffy(EuroSys'22), Faascale(SoCC'24)...]
 - I/O optimization [SAND(ATC'18), SPRIGHT(SIGCOMM'22), RMMMap(EuroSys'24)...]
- Function-centric scheduling:
 - Each function is treated independently [SFS(SC'22), ALPS(ATC'24)...]
 - Workflow-level dependencies and objectives are overlooked
- Gap:
 - Lack of workflow-aware scheduling for end-to-end performance

Problem Definition & Goals

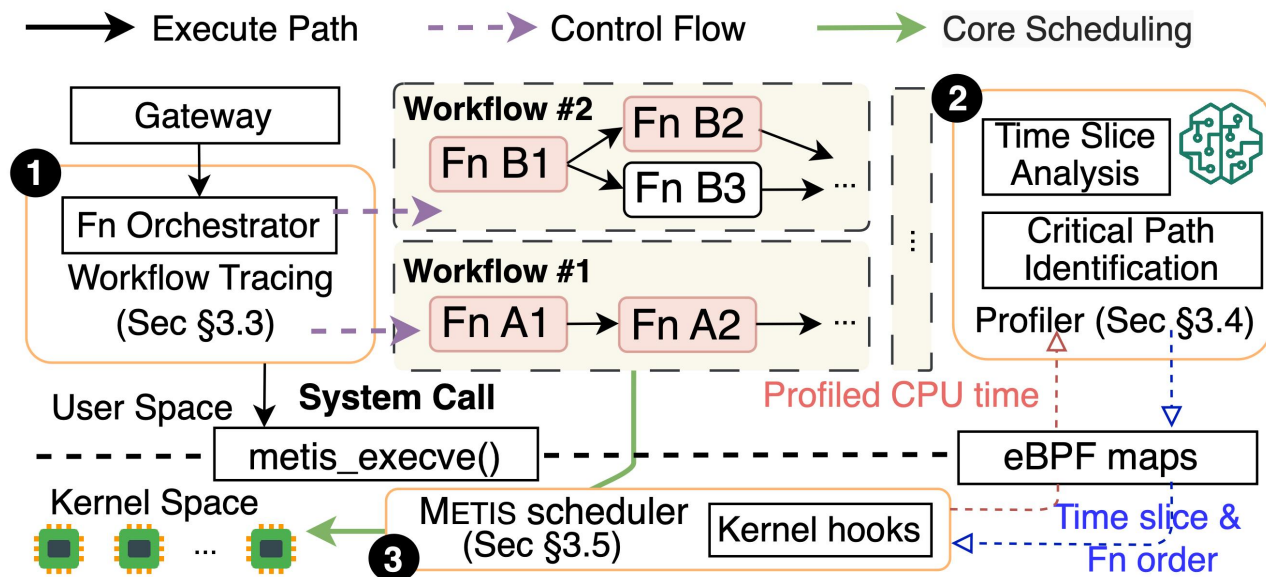
- Problem:
 - OS schedulers lack workflow awareness in serverless environments
 - Critical-path tasks are often blocked, increasing workflow completion time (WCT)
 - Existing solutions require prior knowledge of workflow structures or task durations
- Goals:
 - Enable OS schedulers to optimize workflows as a whole, not just individual functions
 - Minimize end-to-end workflow latency
 - Achieve non-clairvoyant scheduling—no need for prior workflow knowledge

Key Insight: Workflow-Aware Least-Attained Service (WLAS)



METIS Design

Overview of Metis

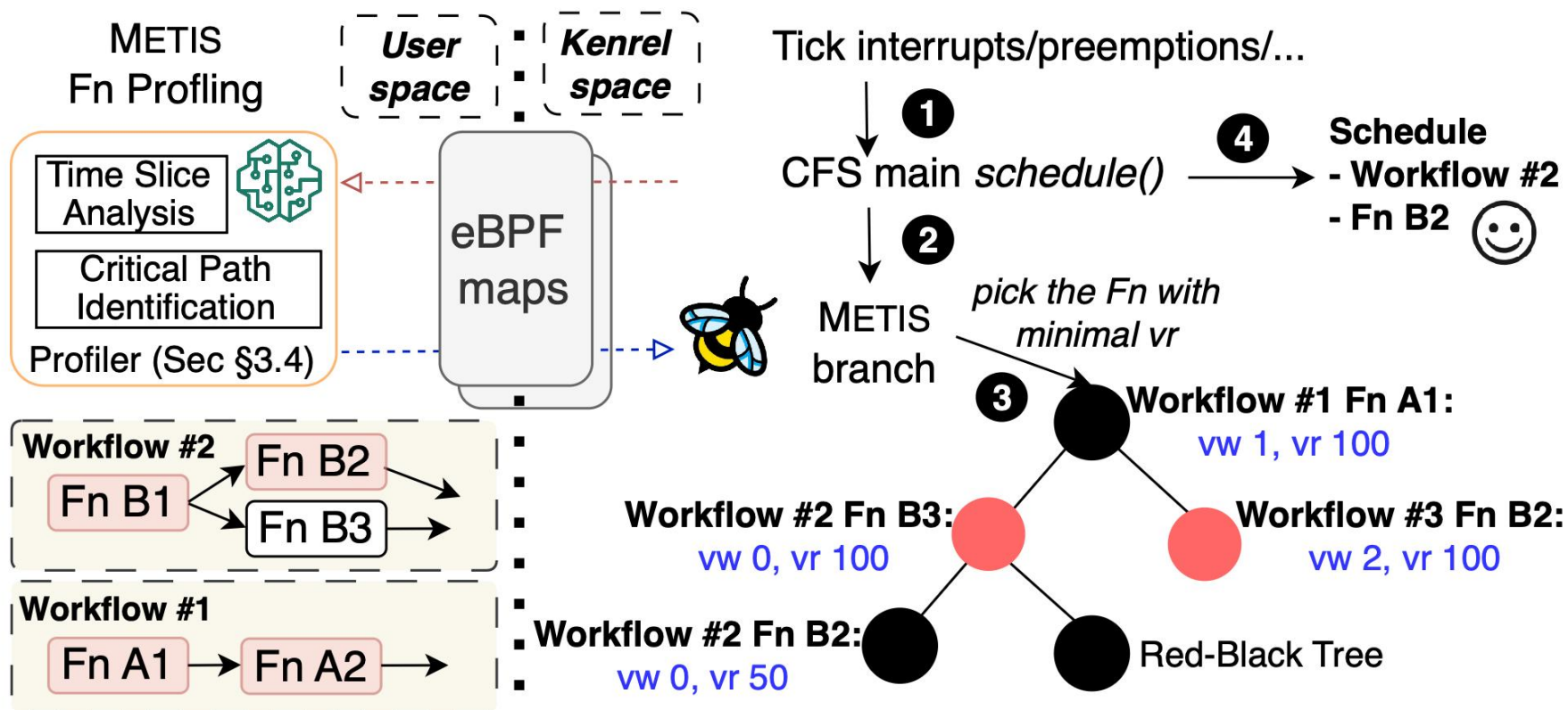


- Three main modules:
 - Workflow tracking—— Identity structure
 - Metric collection—— **critical path estimation**
 - Scheduling decision——ebpf scheduling hints

Key Features of Metis

- Workflow-aware scheduling:
 - Dynamically adapts to workflow structures and dependencies
- Non-clairvoyant design:
 - No need for prior knowledge of workflow structure or task durations
- Low-overhead deployment:
 - Uses eBPF for lightweight, minimally invasive integration with the OS
- Compatibility:
 - Works with mainstream operating systems and serverless platforms

Self-Clocked Workflow-Aware Scheduling



Function scheduling lifecycle in Metis based on workflow-aware prioritization.

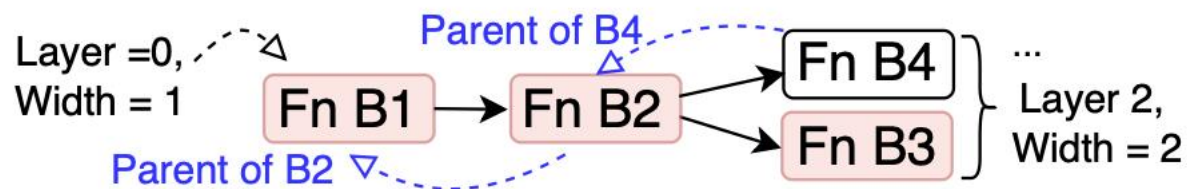
Critical Path Estimation & Function Ordering.

● Critical Path:

- The longest path through the workflow's Directed Acyclic Graph (DAG).
- Determines the minimum possible completion time for the workflow.

● Estimation Method:

- Traverse the DAG to identify the path with the maximum cumulative execution time.



$$CP(w_i) = \max_{\text{path } P \in G_i} \sum_{f \in P} t_f$$

function f 's execution time,

$$p(f_j) = \max_{f_k \in \mathcal{P}(f_j)} \{p(f_k) + t_k\} + \gamma \cdot \text{Width}(L_j)$$

set of parent nodes of f_j

Simulation Results

● Baselines:

- Compared against Linux CFS, FIFO, and ALPS[ATC'24] schedulers

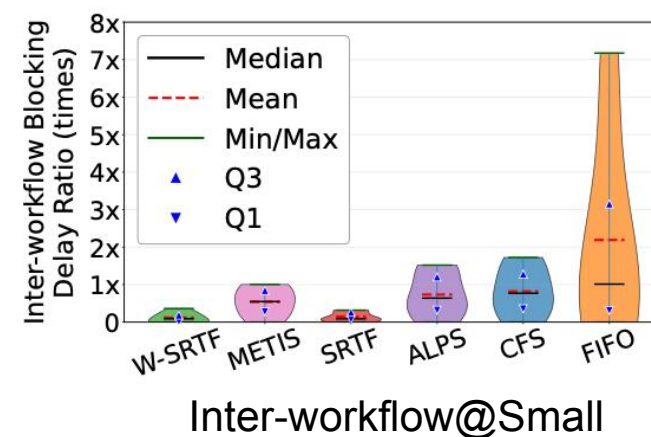
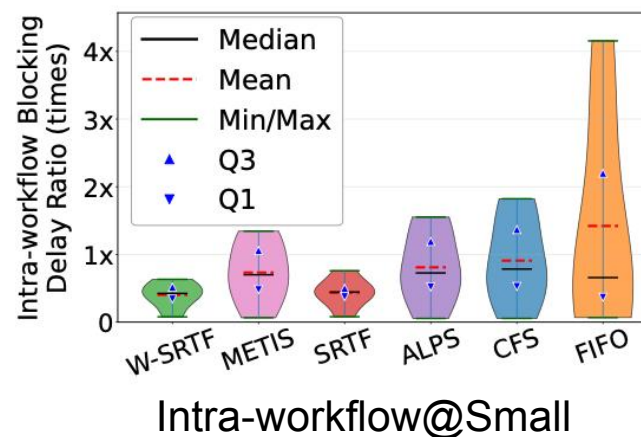
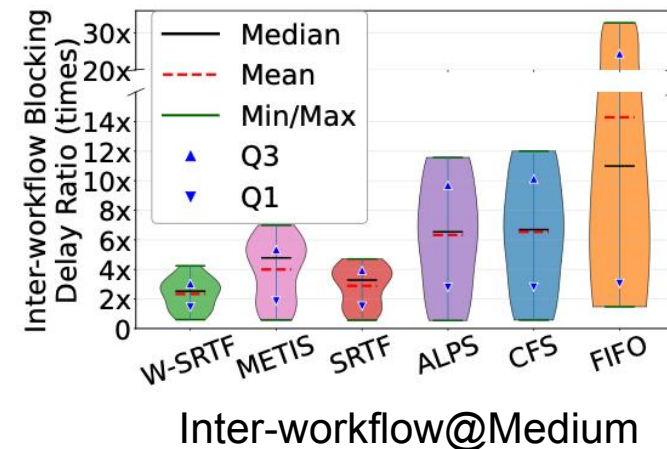
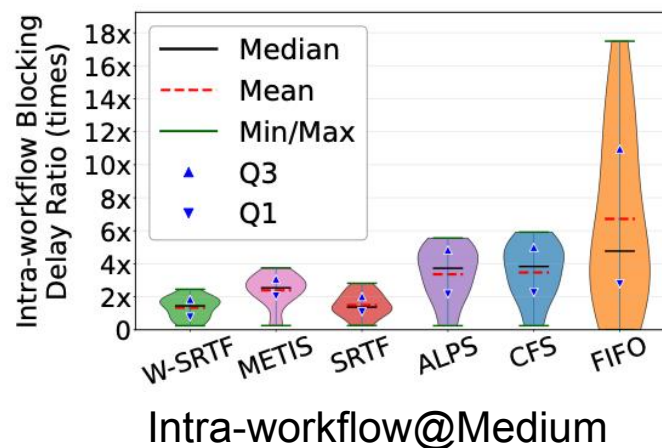
● Evaluation methodology:

- Trace-driven simulations with diverse workflow patterns and scales

● Key results:

- Metis reduces average workflow completion time (WCT) by 31.3%
- 95th percentile latency reduced by 14.3%
- Significant improvement in tail latency (99th percentile).

Synthetic Trace-driven Analysis



End-to-End Workload Performance Comparison

Key characteristics of selected SeBS-Flow benchmarks [EuroSys'25]

● Baselines:

- Compared against Linux CFS, SFS, and ALPS [ATC'24] schedulers

● Experimental Setup:

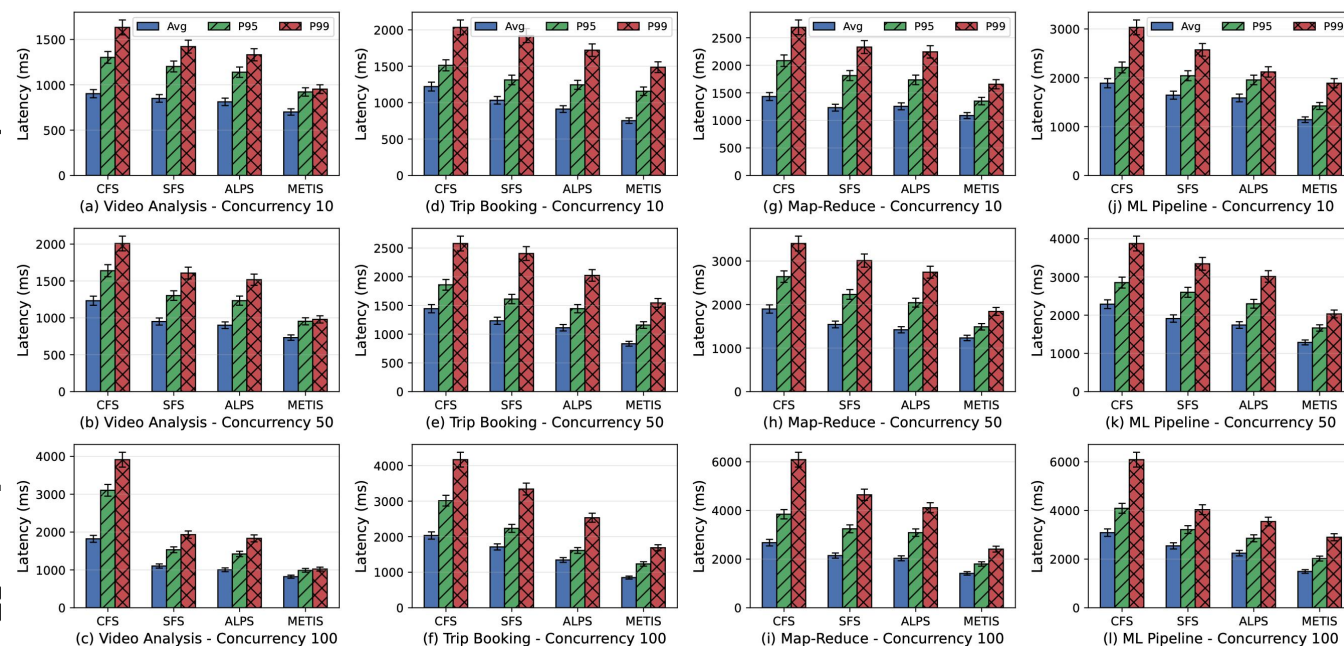
- Four representative workflows (Video Analysis, Trip Booking, MapReduce, ML Pipeline)

- Tested at 10, 50, 100 concurrency levels

● Results:

- Metis reduces average WCT by **31.3%** vs. ALPS
- 95th percentile latency improved by **14.3%**
- 99th percentile latency reduced by up to **73.9%**

Workload	# Func.	Max Par.	DAG Depth	Complexity
Video Analysis	21	16	3	★★
MapReduce	13	10	4	★★★
Trip Booking	8	4	7	★★★★
ML Pipeline	11	8	2	★



Conclusions

- Workflow-centric serverless applications require both fair and efficient OS-level scheduling.
- Metis introduces a non-clairvoyant, workflow-aware scheduler leveraging the novel WLAS algorithm to optimize latency and fairness.
- Comprehensive evaluation demonstrates that Metis consistently surpasses state-of-the-art schedulers.
- Metis achieves up to 58.2% lower workflow completion time and 73.9% better tail latency across diverse workloads.

Thank you!

Any questions?

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